

Before the break:

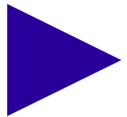
When using Internet technologies, we are confronted with two fundamental questions:

- How to hide *what* is communicated? ◀
- How to hide *who* communicates?

...in the face of an opponent that has *total knowledge of all the IP traffic involved*.

Live example: manually decrypting a ciphertext

"HG"



Live example: manually decrypting a ciphertext

What has just happened?

- The receiver applied a **decryption algorithm**...

"On receiving a message, perform the following two steps:

① *Replace each character by its alphabetical successor: an 'A' becomes a 'B', a 'B' a 'C', etc.*

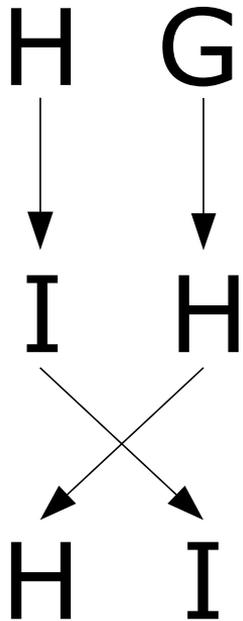
② *Reverse the order: put the last character first, followed by the one-before-last, etc."*

- ...to decrypt a **ciphertext**:

"HG" ① → "IH" ② → "HI"

Live example: manually decrypting a ciphertext

This was an example of a **substitution-permutation cipher**:



Substitution: replace the instances of one symbol by those of another.

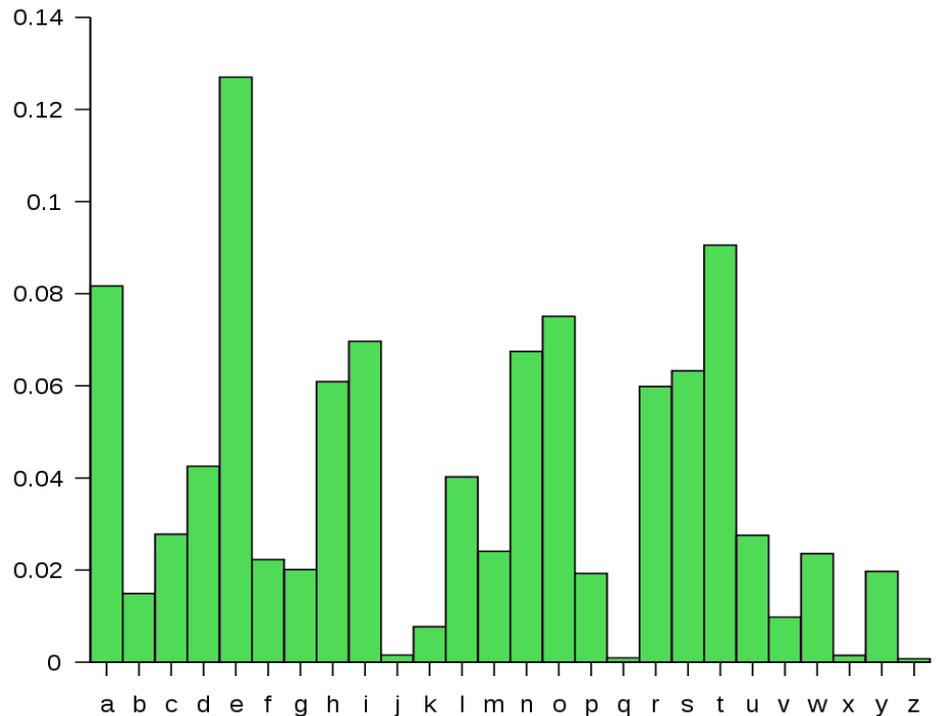
Permutation: change the order in which symbol instances are placed.

Live example: manually decrypting a ciphertext

- The example encryption method uses permutation & substitution to produce an overall ciphertext that is different from the original, **plaintext** message.

- What remains constant however, is that *symbols of the same type* still end up in the ciphertext as *symbols of the same type*.

- This means our ciphertexts will be vulnerable to **letter frequency analysis** →



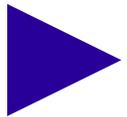
↑ Letter frequencies in English.
Original by Nandhp at Wikimedia Commons.

- To make this attack harder to do, you can group symbol instances together into blocks: a **block cipher**.

Live example: automatic encryption and decryption (using JavaScript)

...However, we will now first *automate* encryption:

- Writing code in JavaScript: the "+x *substitution cipher*".



Live example: automatic encryption and decryption (using JavaScript)

- From the example code, we could further automate the crack () function by adding a dictionary lookup test...

Live example: automatic encryption and decryption (using JavaScript)

- *Everyday intuition:*

“Cracking an encrypted message should be hard.”

- *Now, as code:*

```
ciphertext = encrypt (plaintext, key); // should finish quickly  
plaintext = decrypt (ciphertext, key); // should finish quickly  
plaintext = crack (ciphertext); /* should finish only after the  
                                end of the universe */
```

- ⇒ *So, more precisely:*

The cracking algorithm should consist of many more computational steps than the encryption and decryption algorithms.

Live example: automatic encryption and decryption (using JavaScript)

- In the example, our repeated **decryption** attempts were done using the following code:

```
possible_plaintext = encrypt (ciphertext, -1 * key);
```

- This is an example of **symmetric-key encryption**:

The same key that is used to encrypt the plaintext is also used to decrypt the resulting ciphertext.

- Asymmetric-key encryption is also used, e.g. in public key cryptography:

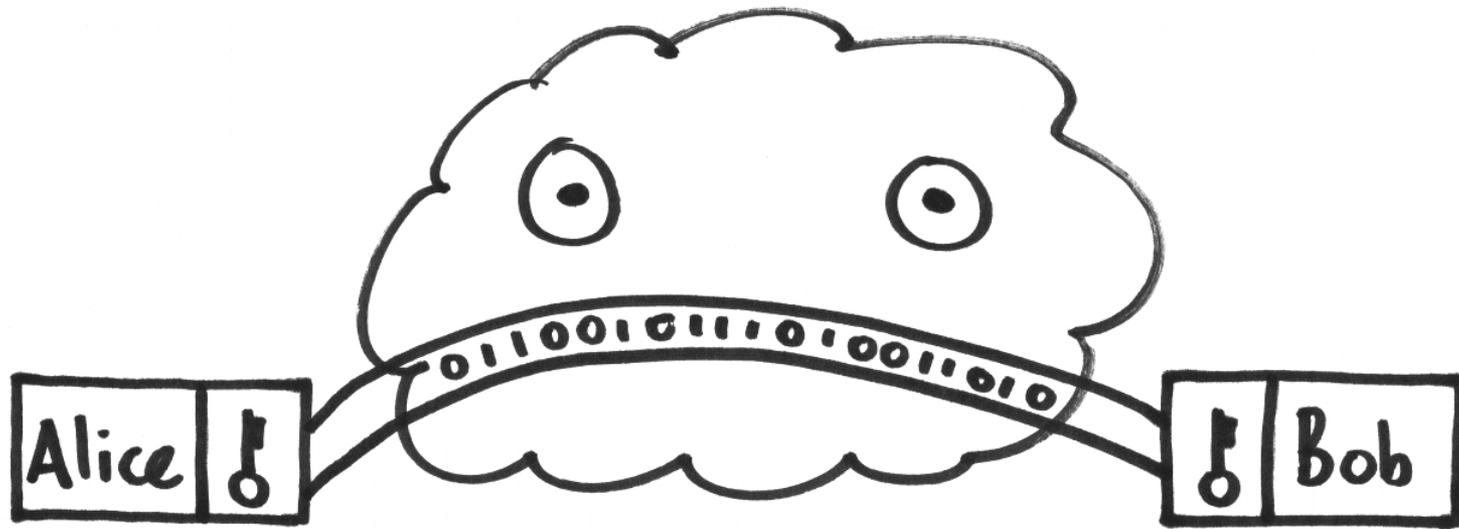
Here, *anybody can encrypt* using a publicly accessible key; but *only the receiver can decrypt* using a second, private key.

Symmetric-key encryption



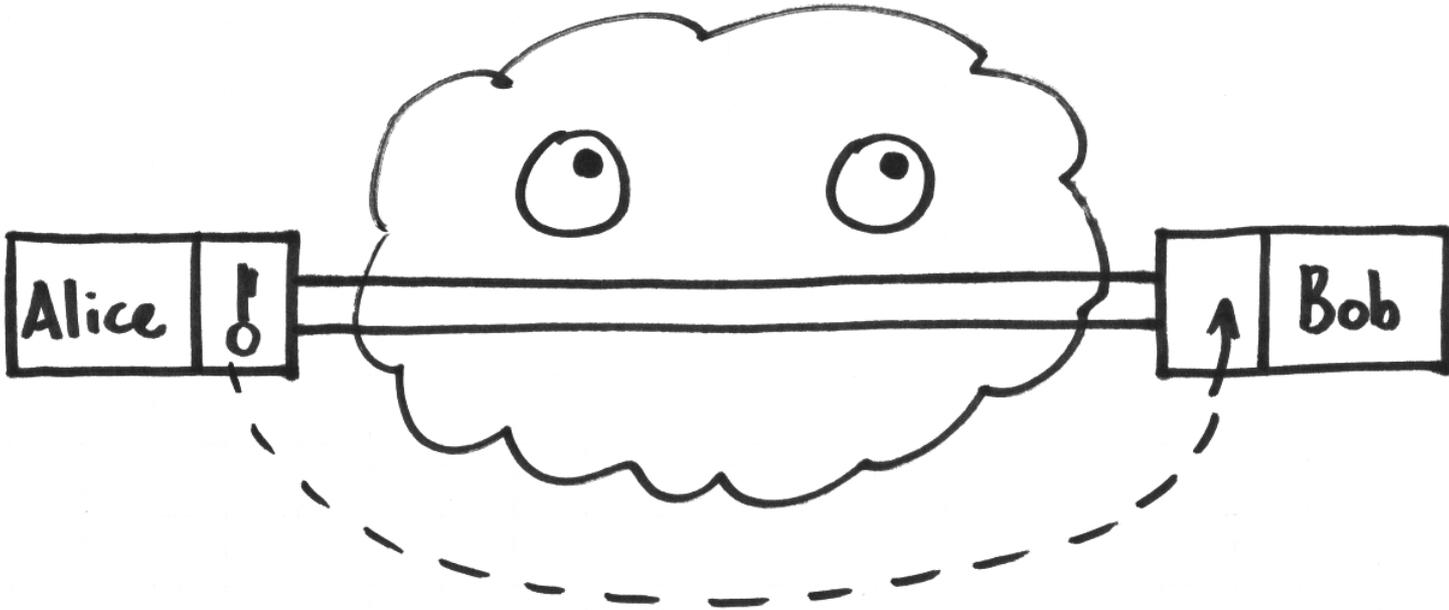
Can be used...

Symmetric-key encryption



...to hide *what* is communicated over the Internet.

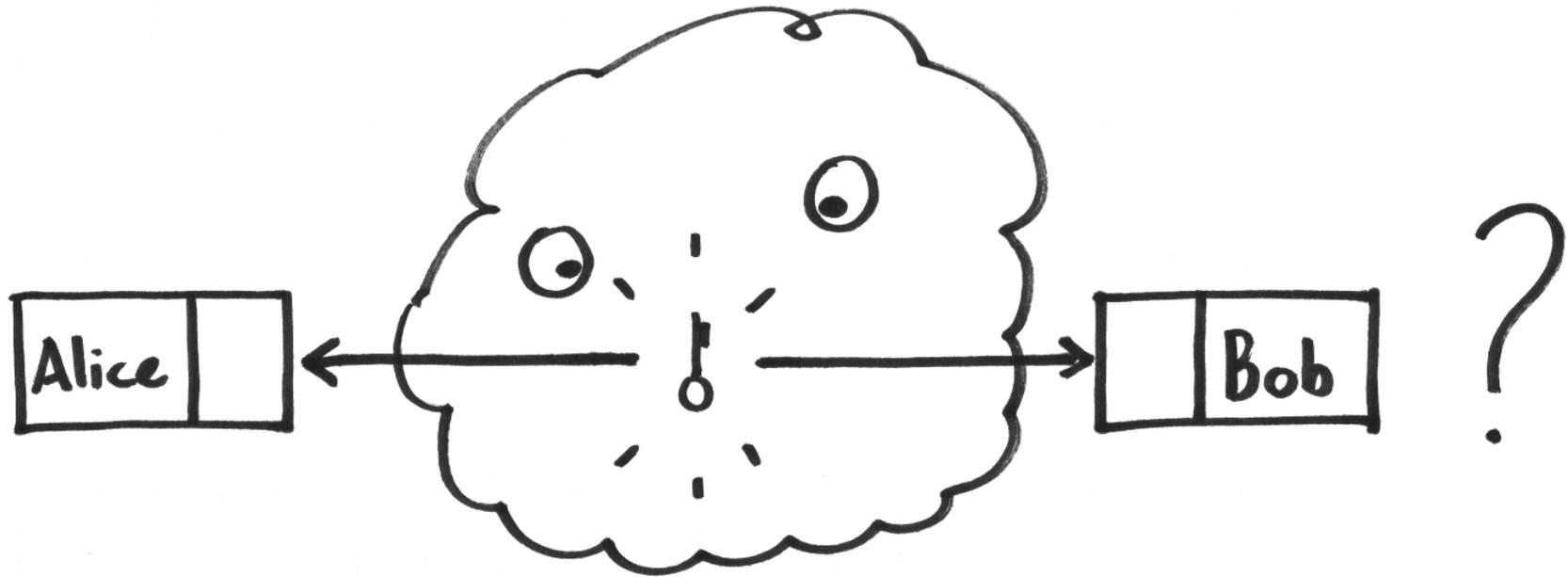
Symmetric-key encryption



- All we need is a covert, secure channel, to exchange the key.

↑ ...Wait a minute!

Symmetric-key encryption



- Remember: our opponent has total knowledge of all IP traffic.

⇒ *Chicken-and-egg problem*: We need a shared key to communicate securely – but we first need to communicate securely, to obtain a shared key. → ...What to do?

Tools to make a secret key with everyone looking

- What to do? We need to grasp three concepts:
 - “An encryption key is a number.”
 - One-way functions.
 - Quasi-commutative functions.

Tools to make a secret key with everyone looking

- A numerical function $f ()$ may combine two input numbers a and b to form one new output number o :

`o = f (a, b); // think of JavaScript here`

- A **commutative** function gives the same output when a and b are swapped.
 - E.g. doing an addition, $(a + b) == (b + a)$ is always true.
 - The general case: " $f (a, b) == f (b, a)$ is always true."
- A **quasi-commutative** function gives the same output in the following argument swap:
 - $f (f (a, b), c) == f (f (a, c), b)$
 - E.g. doing two additions: $(a + b) + c == (a + c) + b$.
 - So: it does not matter if you first combine a with b , then c ; or first with c , then b .

Tools to make a secret key with everyone looking

- A numerical function $f ()$ may combine two input numbers a and b to form one new output number o :

$o = f (a, b); // \text{ think of JavaScript here}$

- For a *one-way* function:
 - When you know a , b , and $f ()$, computing o takes *only a few steps*.
 - But even when you know o , a , and $f ()$, computing b takes *very many steps*.
 - So: computation is quick only in one direction.
 $\Rightarrow f ()$ is “hard to crack” based on its output.

Diffie-Hellman key exchange

- Everyone knows the function $f(\cdot)$ and the constant number g .
- Alice and Bob each secretly pick a random number: a and b .
- Alice privately computes $f(g, a)$, then sends the resulting value to Bob.
- Bob privately computes $f(g, b)$, then sends the resulting value to Alice.
- Everybody can see and record these values!
- But since $f(\cdot)$ is a one-way function, a and b still remain secret.

Diffie-Hellman key exchange

- Bob then uses the **value** of $f(g, a)$ to compute $f(f(g, a), b)$.

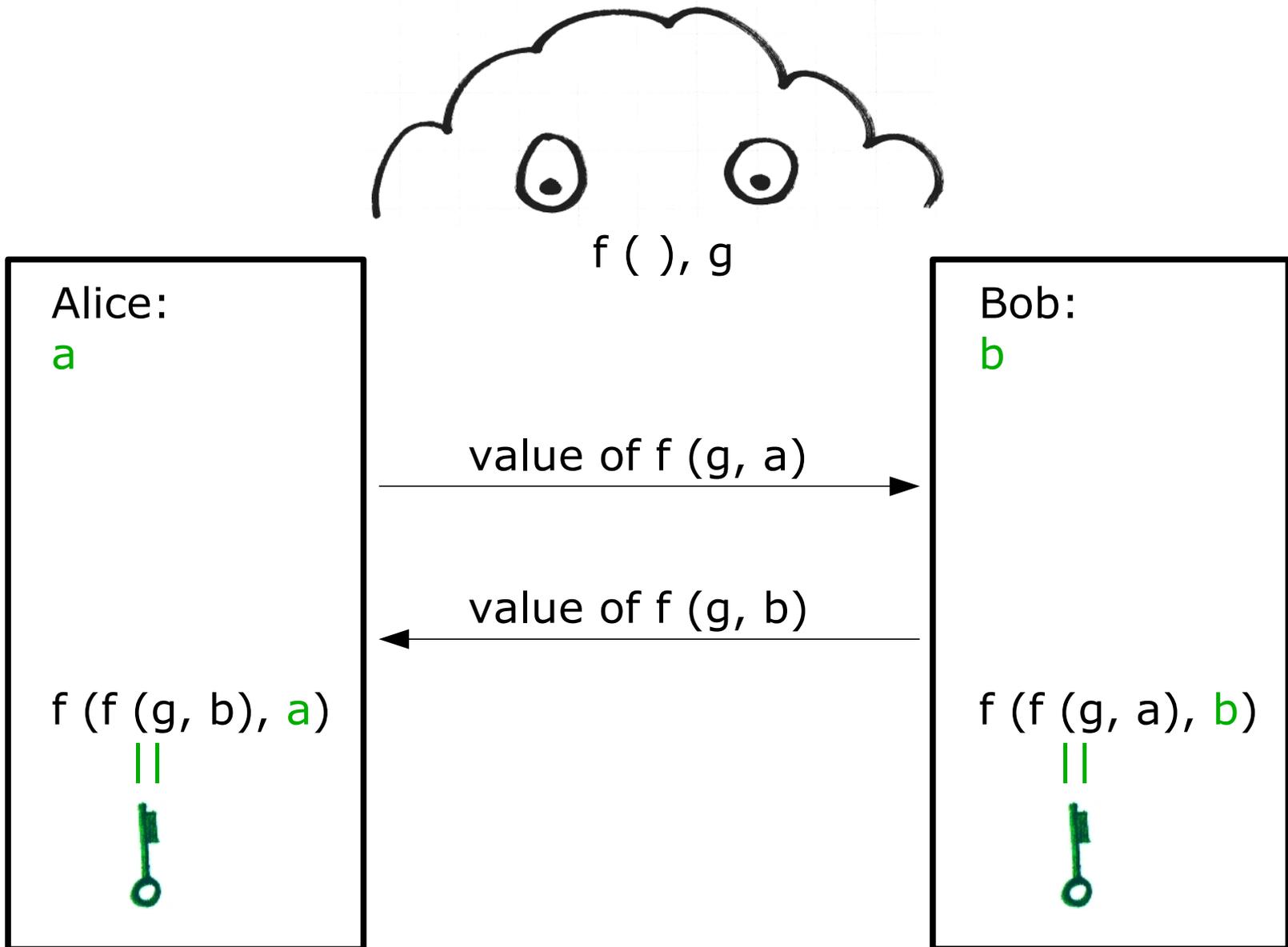
The output is known just to Bob, because only he knows **b**.

- Alice uses the **value** of $f(g, b)$ and computes $f(f(g, b), a)$.

The output is known just to Alice: only she knows **a**.

- But wait – *Bob and Alice actually computed the same value!*
- As $f(\)$ is **quasi-commutative**: $f(f(g, a), b) == f(f(g, b), a)$.
- Bob and Alice now *both* know this number *no one else knows*.
⇒ **They can use it as a symmetric key.**

Diffie-Hellman key exchange, recapitulated



TLS: Transport Layer Security

- *Application layer* protocol which extends transport layer protocols with encryption.
- Diffie-Hellman here is an optional key exchange mechanism.
- TLS has been implemented on top of TCP (but also UDP).
- HTTP Secure (HTTPS): HTTP via TLS (instead of TCP).